Multi-Leader-Follower Games: non cooperative and hierarchical/bilevel interactions

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• Professor in Applied Mathematics at Univ. of Perpignan



Perpignan, France

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- Research topics:

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 - and others....(management of renewable energy plants)

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- Research lab.: PROMES (CNRS)



Optimization /Math. programming



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Hell zone!!!!!

Didier Aussel Multi-Leader-Follower Games: non cooperative and h

An advertisement

A short state of art on Multi-Leader-Follower games, D.A. and A. Svensson, in a book dedicated to Stackelberg, editors A. Zemkoho and S. Dempe, Springer Ed. (2021)



- Lecture 1: Definitions, motivations and well-posedness
- Lecture 2: Some first applications
- Lecture 3: Reformulations (differentiable and non differentiable cases)

Lecture 1

• Definitions of different MLMFG

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Lecture 1

- Definitions of different MLMFG
- Is this well-posed? Is it meaningfull?

Generalized Nash game (GNEP):

$\begin{array}{ c c } \min & & \\ & x_1 \\ & \text{s.t.} \end{array}$	$\theta_1(x_1, x_{-1})$		\min_{x}	$\theta_n(x_n, x_{-n})$
	$\left\{ x_1 \in K_1(x_{-1}) \right.$		s.t.	$\{ x_n \in K_n(x_{-n})$

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So we have n players and they are interacting in a **non cooperative** way (join venture is forbidden or impossible...)

 $ar{x}$ is a Generalized Nash Equilibrium if and only if in case a player i would decide to unilaterally deviate from $ar{x}_i$ (say to $ar{x}_i$) then

"he will loose" := " $\theta_i(\tilde{x}_i, \bar{x}_{-i}) \ge \theta_i(\bar{x}_i, \bar{x}_{-i})$ "!!!!

Motivation for GNEP Power allocation in telecommunication

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Consider a DSL network (Digital Subscriber Line)

- DSL customers connected to the central by dedicaced lines
- wires are bundled together in telephone cables
- electromagnetic coupling \Rightarrow degradation of quality

control variables:

for each wire q and each subcarrier $k,\,p_k^q =$ power allocated for transmission

constraints:

for each wire q: maximum achievable transmission rate R_q (transmission quality)

It depends of $(p_k)_{k=1,N}$ (power allocations across available subcarriers for q) and of $p^{-q} = (p^r)_{r \neq q}$ (strategies of the other wires)

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Control variables:

$$R_q(p^q, p^{-q}) = \sum_{k=1}^N log(1 + sinr_k^q)$$

where

$$sinr_{k}^{q} = \frac{|H_{k}^{qq}|^{2}.p_{k}}{\sigma_{q}^{2}}(k) + \sum_{r \neq q} |H_{k}^{qr}|.p_{k}^{r}.$$

(Signal-t-Interference Noise Ratio) Garantee of minimal transmission rate $R_q(p^q, p^{-q}) \ge R_q^*$.

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Model:

each wire wants is a player of the game whose objective is to minimize to total power used for transmission, with the constraint that the maximum transmission rate is at least R_a^* , that is

for any
$$q$$
, solve (P_q) min $\sum_k = 1^k p_k^q$
 $s.t. \begin{cases} R_q(p^q, p^{-q}) \ge R_q^* \\ p_k^q \ge 0 \end{cases}$

Theorem

Assume that

- for every player ν, the loss function θ_ν is continuous on ℝⁿ and semistrictly quasiconvex with respect to the ν-th variable;
- the set X is nonempty, convex and compact.

Then the Nash equilibrium problem admits a solution.

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Theorem (Debreu)

Assume that

- for every player ν, the loss function θ_ν is continuous on ℝⁿ and quasiconvex with respect to the ν-th variable;
- the set-valued map is continuous with nonempty convex compact values
- . Then the generalized Nash equilibrium problem admits a solution.

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Bilevel optimisation

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Bilevel problem:

$$\begin{array}{ll} \min_{x} & \theta(x,y) \\ \text{s.t.} & \begin{cases} x \in X(y) \\ \min_{y} & \phi(x,y) \\ \text{s.t.} & y \in Y(x) \end{cases}$$

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• Ambiguities...

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• Ambiguities...

• Leader or Follower points of view?

BL = Single-Leader-Single-Follower game

A Bilevel Problem consists in an **upper-level/leader's** problemthis

where $\emptyset \neq X \subset \mathbb{R}^n$ and S(x) stands for the solution set of its lower-level/follower's problem

$$\min_{y \in \mathbb{R}^m} \quad f(x, y) \\ s.t \quad g(x, y) \le 0$$

A trivial example

Consider the following simple bilevel problem

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$$\min_{x \in \mathbb{R}} x \quad x \in [-1, 1]$$

s.t.
$$\begin{cases} x \in [-1, 1] \\ y \in S(x) \end{cases}$$

with S(x) = "y solving

$$\min_{y \in \mathbb{R}} \quad -xy \\ s.t \quad x^2(y^2 - 1) \le 0 "$$

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Lower level problem:

$$\min_{y \in \mathbb{R}} \quad -x.y \\ s.t \quad x^2(y^2 - 1) \le 0$$

Note that the solution map of this convex problem is

$$S(x) := \begin{cases} \{1\} & x < 0\\ \{-1\} & x > 0\\ \mathbb{R} & x = 0 \end{cases}$$

Thus for each $x \neq 0$ there is a unique associated solution of the lower level problem

Lower level problem:

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An Optimistic Bilevel Problem consists in an upper-level/leader's problem

$$\min_{x \in \mathbb{R}^n} \min_{y \in \mathbb{R}^m} \quad F(x, y) \\ s.t. \quad \begin{cases} x \in X(y) \\ y \in S(x) \end{cases}$$

where $\emptyset \neq X \subset \mathbb{R}^n$ and S(x) stands for the solution set of its lower-level/follower's problem

$$\min_{y \in \mathbb{R}^m} \quad f(x, y) \\ s.t \quad g(x, y) \le 0$$

An *Pessimistic Bilevel Problem* consists in an **upper-level/leader's problem**

$$\min_{x \in \mathbb{R}^n} \max_{y \in \mathbb{R}^m} \quad F(x, y) \\ s.t. \quad \begin{cases} x \in X(y) \\ y \in S(x) \end{cases}$$

where $\emptyset \neq X \subset \mathbb{R}^n$ and S(x) stands for the solution set of its lower-level/follower's problem

$$\min_{y \in \mathbb{R}^m} \quad f(x, y) \\ s.t \quad g(x, y) \le 0$$

Ambiguity: the most simple

And of course the "confortable situation" corresponds to the case of a unique response

$$\forall x \in X, \quad S(x) = \{y(x)\}.$$

Then

$$\min_{x \in \mathbb{R}^n} \quad F(x, \mathbf{y}(x)) \\ s.t. \quad \left\{ \begin{array}{l} x \in X(\mathbf{y}(x)) \end{array} \right.$$

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For example when

for any x, $g(x, \cdot)$ is quasiconvex and $f(x, \cdot)$ is strictly convex.

An *"Selection-type"* Bilevel Problem consists in an upper-level/leader's problem

$$\begin{array}{l} \min_{x \in \mathbb{R}^n} & F(x, y(x)) \\ s.t. & \begin{cases} x \in X(y(x)) \\ y(x) \text{ is a uniquely determined selection of } S(x) \end{cases}$$

J. Escobar & A. Jofré, Equilibrium Analysis of Electricity Auctions (2011)

Recently, D.Salas and A. Svensson proposed a **probabilistic approach**:

- Consider a probability on the different possible follower's reactions
- Minimize the expectation of the leader(s)

Instead of considering the previous (optimistic) formulation of BL:

$$\min_{x \in \mathbb{R}^n} \min_{y \in \mathbb{R}^m} \quad F(x, y) \\ s.t. \quad \begin{cases} x \in X \\ y \in S(x) \end{cases}$$

Instead of considering the previous (optimistic) formulation of BL:

$$\min_{x \in \mathbb{R}^n} \min_{y \in \mathbb{R}^m} \quad F(x, y) \\ s.t. \quad \begin{cases} x \in X \\ y \in S(x) \end{cases}$$

one can define the (optimistic) value function

$$\varphi_{\min}(x) = \min_{y} \{ F(x, y) : g(x, y) \le 0 \}$$

$$\tag{1}$$

and the BL problem becomes

 $\begin{array}{ll} \min_{x \in \mathbb{R}^n} & \varphi_{min}(x) \\ s.t. & x \in X \end{array}$

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Instead of considering the previous (pessimistic) formulation of BL:

$$\min_{x \in \mathbb{R}^n} \max_{y \in \mathbb{R}^m} \quad F(x, y) \\ s.t. \quad \begin{cases} x \in X \\ y \in S(x) \end{cases}$$

one can define the (pessimistic) value function

$$\varphi_{max}(x) = \max_{y} \{ F(x, y) : g(x, y) \le 0 \}$$
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$$\min_{x \in \mathbb{R}^n} \min / \max_{y \in \mathbb{R}^m} F(x, y) s.t. \begin{cases} x \in X & \mathbf{vs} & \min_{x \in \mathbb{R}^n} \varphi_{\min/max}(x) \\ y \in S(x) & s.t. & x \in X \end{cases}$$

It immediately raises the question

What is a solution??

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It immediately raises the question

What is a solution??

- an optimal x = leader's optimal strategy?
- an optimal couple (x, y) = couple of strategies of leader and follower?